

Compiling and Linking C / C++ Programs

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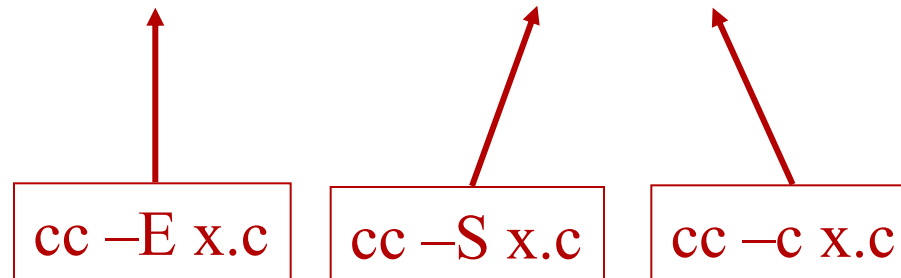
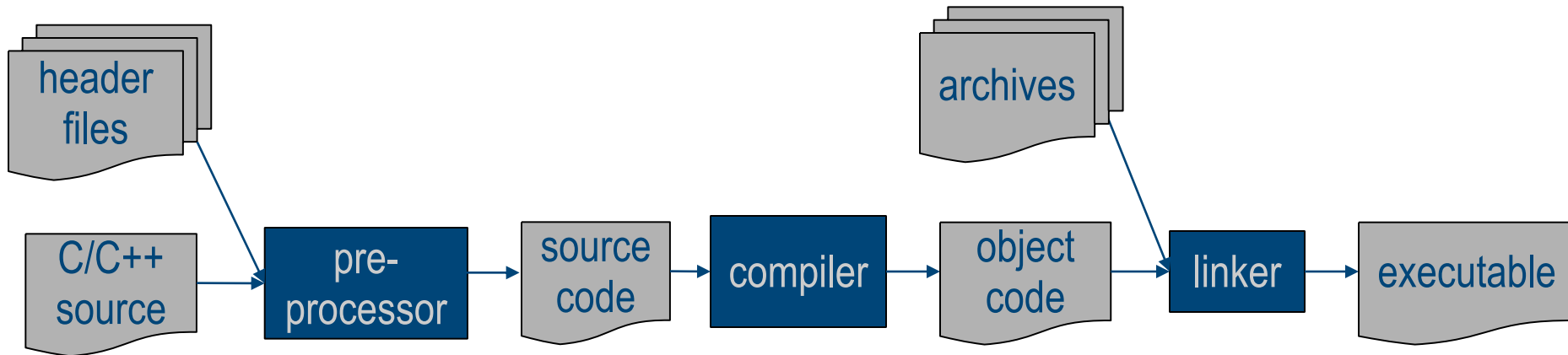
office: room 88, Research 1



CPU @ Work

- ...watch your code like you never have seen it before!

Compile/Link Steps Overview



File Extension Conventions

- C source code .c
- C include file .h
- C++ source file .cc , .C, .cxx, .c++, .cp, .cpp
- C++ header file .hh, .hpp
- Object file (relocatable) .o
- Executable no extension (Windows: .com, .exe)
- Library
 - static .a
 - dynamic .so

The C Preprocessor

- Purpose:
 - Define commonly used constants, code fragments, etc.
 - Conditional compilation (code activation depending on external settings)
- Main mechanism: replace by **textual substitution**
 - No idea about semantics (parentheses, semicolons, ...) !!
 - Does not follow C syntax
- Preprocessor directives
 - `#include`
 - `#define`
 - `#if / #ifdef`
 - ...plus more

```
#define X 1
```

```
const int x = 1;
```

Using Preprocessor Directives

- Conditional compilation

- Include guard in header files, eg in `mystdio.h`:

```
#ifndef _MYSTDIO_H_
#define EOF (-1)
#define NULL 0L
#define _MYSTDIO_H_
#endif _MYSTDIO_H_
```

- Include files

- `#include <stdio.h>` – taken from predefined location
- `#include "myclass.h"` – taken from local directories

- Where to find include files?

- Standard locations: `/usr/include, /usr/local/include, ...`
- Specified locations `cc -I/home/project/include`

- Can also pass definitions

- `cc -DCOMPILE_DATE=\" `date` \" -DDEBUG`


Common Preprocessor Pitfall

- Use parentheses!!!

- bad:

```
#define mult(a,b) a*b
main()
{
printf( "(2+3)*4=%d\n", mult(2+3,4) );
}
```

```
printf( "(2+3)*4=%d\n", 2+3*4 );
```



- good:

```
#define mult(a,b) ((a)*(b))
main()
{
printf( "(2+3)*4=%d\n", mult(2+3,4) );
}
```

```
printf( "(2+3)*4=%d\n", ((2+3)*(4)) );
```



The C(++) Compiler

- Task: Generate (relocatable) **machine** („object“) **code** from source code
- **Relocation**: code can sit at different places in address space
- Address space classified into „segments“
 - Code, text, data, ...
- Note: OS (with HW support) uses this to implement **user address space**
 - Actual main memory address = base address + relative address
 - Base address kept in segment register, added dynamically by CPU
 - Security: program cannot access base register ("privileged mode"), hence cannot address beyond its segment limits

Object Files

- Contain code for a program fragment (module)
 - Machine code, constants, size of static data segments, ...

```

$ nm rserver_main.o
0000000c D clientTimeout
          U __cxa_allocate_exception
          U __cxa_begin_catch
          U __cxa_end_catch
          U __cxa_free_exception
          U __cxa_throw
00000004 B debugOutput
          U free
          U getenv
00000120 B globalHTTPPort
00000000 T main
          U memset

```

or objdump

External Functions & Variables

- Module **server**:
Variable **sema** allocated in data segment
- Module **client**:
functions obtain **sema** address by
 - Module **server** offset
+ local address **sema**
- Cross-module addressing rules:
 - (no modifier) = locally allocated, globally accessible
 - **static** = locally allocated, locally accessible
 - **extern** = allocated in other compilation unit
- Why is this wrong?
 - **extern int sema = 1;**

```
int sema = 0;

int serverBlock()
{
  if (sema==0)
    sema = 1;
  return sema;
}
```

```
extern int sema;

int clientBlock()
{
  if (sema==0)
    sema = 1;
  return sema;
}
```

Name Mangling

- Problem: classes convey complex naming, not foreseen in classic linkage
 - Classes, overloading, name spaces, ...
 - Ex: `MyClass1::myFunc()`
`MyClass2::myFunc()`
 - But only named objects in files, flat namespace
- Solution: **name mangling**
 - Compiler **modifies names** to make them unique (prefix/suffix)
 - Ex: `Transaction::begin()`
 → `_ZN13r_Transaction5beginENS_8r_TAModeE`
- Every compiler has its individual mangling algorithm!
 - Code compiled with different compilers is **incompatible**

Name Mangling (contd)

- Q: My linker cannot find function `flatFunc()`, although I link against library `libff.a` which definitely contains `flatFunc()`.
 - Proof: `nm libff.a | grep flatFunc`
- A: You're probably linking with C code.
Tell the C++ compiler that it's C, not C++, to avoid name mangling:

```
extern "C"
{
    void flatFunc();
}
```

The Linker/Loader

- Task: generate *one* executable file from *several* object and library files
 - Read object files
 - Resolve addresses from (relocatable) code
 - Link runtime code (start/exit handling!)
 - Add all code needed from libraries
 - If required: establish stubs to dynamic libraries
 - Write executable code into file, set magic number, etc.

- cc, g++, etc. have complex logics inside
 - can silently *invoke* linker, don't link themselves!
 - Common shorthand: `cc -o x x.c`

- Ex: `ld -o x /lib/crt0.o x.o -lc`

*John R. Levine:
Linkers and Loaders.
Morgan Kaufmann, 1999*

What It Really Looks Like

```
if (DEBUG_LEVEL >= DEBUG_MEM) {  
    memrec_add_var(&malloc_rec, filename, line, temp, size);  
}  
return (temp); void *ptr, size_  
}  
  
realloc(const char *var, const char *filename, unsigned long line, void *ptr, size_t size)  
{  
    void *temp;  
  
#ifdef MALLOC_CALL_DEBUG  
    ++realloc_count;  
    if (!(realloc_count % REALLOC_MOD)) {  
        D_MEM(("Calls to realloc(): %d\n", realloc_count));  
    }  
#endif  
    line, size_t count,  
    D_MEM(("Variable %s (%8p -> %lu) at %s:%lu\n", var, ptr, (unsigned long) size, filename, line));  
    if (ptr == NULL) {  
        temp = (void *) libast_malloc(__FILE__, __LINE__, size);  
    }  
    else {  
        temp = (void *) realloc(ptr, size);  
        ASSERT_RVAL(temp != NULL, ptr);  
        if (DEBUG_LEVEL >= DEBUG_MEM) {  
            memrec_chg_var(&malloc_rec, var, filename, line, ptr, temp, size);  
            memrec_add_var(&malloc_rec, filename, line, temp, size);  
        }  
    }  
    return (temp);  
}  
  
const char *filename, unsigned long line, size_t count, size_t size)
```

```
CC -V X.C
```

```
cc1 ... x.c ... -o /tmp/ccWs4dqa.s  
as ... -o /tmp/cckBDoD2.o /tmp/ccWs4dqa.s  
collect2 ... -o x /lib/ld-linux.so.2 \\  
    crt1.o crti.o crtbegin.o /tmp/cckBDoD2.o  
    -lgcc -lgcc_eh -lc -lgcc -lgcc_eh  
    crtend.o crtn.o
```

Strip

- By default, executable contains symbol tables
 - Function names, addresses, parametrization
 - Static variables
 - ...some more stuff

- Disadvantages:
 - Allows reverse engineering (gdb!)
 - Substantially larger code files

- Before shipping: strip executables
 - **file rasserver**
`rasserver: ELF 32-bit LSB executable, Intel 80386, version 1 (SYSV), for GNU/Linux 2.2.5, dynamically linked (uses shared libs), not stripped`
 - **strip rasserver**

Libraries (Archive Files)

- Library = archive file containing a collection of object files
 - Code fragments (classes, modules, ...)
 - `ar rv libxxx.a file1.o file2.o ...`

- Object files vs. Libraries
 - Object file linked in **completely**, from library **only what is actually needed**

- Static vs. Dynamic
 - **Static library**: code is added to the executable, just like object file; not needed after linkage
 - **Dynamic library**: only stub linked in, runtime system loads; needed at runtime (version!)

- Naming conventions (Linux)
 - Static libraries: **libxxx.a**
 - Dynamic libraries: **libxxx.so**
 - link with: `ld ... -lxxx`

Dynamic Libraries

- How to **find** my dynamic libraries?
 - `LD_LIBRARY_PATH` variable, similar to `PATH`: set before program start

- How to **know about use** of dynamic libraries?
 - `$ ldd rserver`

```

linux-gate.so.1 => (0xffffe000)
libstdc++.so.5 => /usr/lib/libstdc++.so.5 (0x40028000)
libm.so.6 => /lib/tls/libm.so.6 (0x400e5000)
libgcc_s.so.1 => /lib/libgcc_s.so.1 (0x40128000)
libc.so.6 => /lib/tls/libc.so.6 (0x40130000)
libresolv.so.2 => /lib/libresolv.so.2 (0x4029c000)
          
```

Schematic Program Run

- OS:
 - Open file
 - Look at first page:
magic number, segment sizes, etc.
 - Allocate segments
(code, runtime stack, heap, ...)
 - Read code file into code segment
 - Set up process descriptor
(external resources, limits, ...)
 - Pass control to this process
 - Handle system calls
 - Terminate program,
free process slot and resources
- Application program:
 - Set up runtime environment
(`argv/argc`, ...)
 - Call `main()`
 - On system calls, interrupts, etc.:
pass control to OS
 - Upon `exit()`,
or `main()`'s `return`,
or a forced abort:
clean up (close file descriptors, sockets, ...),
pass back to OS

Summary

- To create executable program, you must perform:
 - Preprocess – textually expands definitions, condition-guarded code pieces
 - Compile – translates source code into relocatable machine code („object code“)
 - Link – bind object files and archives into executable program

```
cc -o x x.c
```

=

```
cpp x.c x.cpp
cc -o x.o -c x.cpp
ld -o x /lib/crt0.o x.o -lc
```

Summary (contd.)

- A word about code quality

- Set compiler to max rigidity: `cc -W -Wall ...`
- Eliminate *all* warnings

- Finally, the formatting war:

ANSI C:

```
void foo()
{
    myAction();
}
```

Kernighan/Ritchie:

```
void foo() {
    myAction();
}
```

- The answer: whatever style, **use one coherently**

- Use automatic beautifier (see my SE course page for some)